https://pingo.scc.kit.edu/815630
Document Listing

- similar to document retrieval (next lecture)
- get all documents containing a phrase

Definition: Document Listing

Given a collection of $D$ documents $\mathcal{D} = \{d_1, d_2, \ldots, d_D\}$ containing symbols from an alphabet $\Sigma = [1, \sigma]$ and a pattern $P \in \Sigma^*$, return all $j \in [1, D]$, such that $d_j$ contains $P$.

- $d_1 = ATA$
- $d_2 = TAAA$
- $d_3 = TATA$

And for queries:
- $P = TA$ is contained in $d_1$, $d_2$, and $d_3$
- $P = ATA$ is contained in $d_1$ and $d_3$
Definition: Document Concatenation

Given a collection of \( D \) documents \( D = \{d_1, d_2, \ldots, d_D\} \) containing symbols from an alphabet \( \Sigma = [1, \sigma] \) where each document ends with a special symbol \( \# \), the string

\[ C = d_1 d_2 \ldots d_D \# \]

is called the concatenation of the documents with \( \# \notin \Sigma \) and \( \$ < \# < \alpha \) for all \( \alpha \in \Sigma \)

\[ N = |C| = \sum_{i=1}^{D} |d_i| \]
**Suffix Array for Document Concatenation**

- Given a document concatenation $C$, build the suffix array.
- Requires $O(n)$ time.
- Entries in suffix array correspond to documents.

**Definition: Document Array**

Given a document concatenation $C$ and its suffix array $SA$, the **document array** $DA$ is defined as

$$\text{DA}[i] = j \iff \sum_{k=1}^{j-1} |d_k| < SA[i] \leq \sum_{k=1}^{j} |d_k|$$

for $i > 1$ and $\text{DA}[1] = 0$.

---

**PINGO** are the document and suffix array sufficient to return all documents?
Naive Document Listing

- given document concatenation $C$, its suffix array $SA$, and document array $DA$
- enhance suffix array to do pattern matching in $O(|P|)$ time only briefly discussed in lecture
- find interval in suffix array matching $P$
- report all documents in interval in $DA$
- problem: $O(|P| + N)$ query time very bad

- is there a better solution?
- better query time
- better (or at least equal) space requirements?

$P = TA$
Definition: Chain Array

Given document concatenation $C$, its suffix array $SA$, and document array $DA$, the chain array $CA$ is defined as

$$CA[i] = \max\{j < i : DA[j] = DA[i]\} \cup \{0\}$$

- chains same documents together
- find lexicographically smaller suffix of same document
- use it to report documents just once
- build RMQ data structure for $CA$

$P = TA$

PINGO is the chain array with RMQs enough to list all documents in optimal time?
Optimal Time Document Listing (2/2)

- given document concatenation $C$, its suffix array $SA$, document array $DA$, and chain array $CA$ with RMQ data structure
- find interval $SA[s, e]$ as before
- report document $DA[m]$ only if $CA[m] < s$ for $m \in [s, e]$

- find all positions where $CA[m] < s$ with RMQs
- get arg min of $CA$ in interval and report $DA[m]$ if $CA[m] < s$
- split interval in $[s, m - 1]$ and $[m + 1, e]$ and recurse
- ignore intervals where nothing is reported

Lemma: Optimal Document Listing

Listing all documents containing a pattern $P$ can be done in $O(|P| + occ)$ time
Top-$k$ Document Retrieval for Single-Term Frequencies

Definition: Top-$k$ Document Retrieval

Given a collection of $D$ documents $\mathcal{D} = \{d_1, d_2, \ldots, d_D\}$ containing symbols from an alphabet $\Sigma = [1, \sigma]$, a pattern $P \in \Sigma^*$, and a threshold $k$, return the top-$k$ documents $j \in [1, D]$, such that $d_j$ contains $P$ most often.

- retrieve $occ$ distinct documents where $P$ occurs
- determine frequency of $P$ in each document
- maintain min-heap of (frequency,document)-pairs of size $k$
- total time: $O(|P| + occ(\lg k + \lg N))$

occ can be $N$
- can we do better

optimal solution: $O(|P| + k)$ query time in $O(N \lg N)$ bits [NN12]
- now: $O(|P| + k \lg N)$ [GKN17]
Recap: Suffix Tree

Definition: Suffix Tree [Wei73]
A suffix tree (ST) for a text $T$ of length $n$ is a compact trie over $S = \{ T[1..n], T[2..n], \ldots, T[n..n] \}$ suffixes are prefix-free due to sentinel

Let $G = (V, E)$ be a compact trie with root $r$ and a node $v \in V$, then
- $\lambda(v)$ is the concatenation of labels from $r$ to $v$
- $d(v) = |\lambda(v)|$ is the string-depth of $v$

string depth $\neq$ depth
a generalized suffix tree is a suffix tree for a set of strings

document concatenation is a set of strings

Mark Document Numbers

- mark all leaves with $DA$-entry $i$
- add $i$ to nodes that are lowest common ancestor of two leaves marked with $i$
Inner Node Names

- leaf index is rank of suffix in $[1, N]$ in leaf
- each inner node gets $v$ gets $id(v)$, which is the leaf index of rightmost leaf in leftmost child

- $id(v) \neq id(w)$ for all inner nodes $v \neq w$
- $id(v) \in [1, N]$ 
- $id(v) - 1 \in [lb(v), rb(v)]$, with $[lb(v), rb(v)]$ being $v$'s suffix array interval

example on the board 📈
connect node with id $i$ to closest ancestor containing id $i$

- nodes marked with id $i$ correspond to suffix tree of $d_i$
- document id $i$ occurs at most $|d_i|$ times in leaves and $|d_i| - 1$ times in inner nodes
- there are at most $O(N)$ document ids in the generalized suffix tree
to retrieve documents containing pattern $P$
select locus of $P$ 1 first node $v$ with $P$ is prefix of $\lambda(v)$

per document at most one pointer leaves subtree of locus $v$
associate each pointer with number of occurrences of documents in pointers source (weight)
pointer of document $i$ leaving subtree has maximum weight of all document $i$ pointers in subtree
document listing is listing all documents of pointers leaving subtree
Representing Pointers on a Grid (1/2)

- now: report top-\(k\) documents
- represent pointers in a grid
- for simplicity only weights \(\geq 2\) starting at inner node

- assign each pointer to \((x, y)\)-coordinate
  - \(x\): \text{id}(source)
  - \(y\): \text{d}(target)
- each point is associated with pointers weight

- given a locus \(v\), all pointers leaving the subtree have \(y\)-coordinate < \(\text{d}(v)\)
Representing Pointers on a Grid (2/2)

- grid can be represented using wavelet tree
- range maximum query for each level

Answering Queries

- find string depth of locus in suffix tree
- answer range query in grid
- if represented as wavelet tree, use RMQs on each level to report top-$k$ documents
- if $\leq k$ documents, use document listing
- total time: $O(m + k \lg N)$

PINGO how can we represent the pointers in a grid?

example range queries in wavelet trees on the board
Conclusion and Outlook

This Lecture
- document listing
- top-$k$ document retrieval (single term frequency)

Next Lecture
- inverted indices (by Tim Niklas Uhl)
Oral Exam

- registration is open
- oral exams 24.02. and 20.03.
- alternative dates possible
- register with our secretary
Bibliography I


