Lecture 3: Construction of Suffix Arrays

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Fig. 2. Taxonomy of suffix array construction algorithms.

Normally prefix-doubling algorithms initialize $SA_1$ for $h = 1$ using a linear-time bucket sort. The main idea [Karp et al. 1972] is as follows:

**Observation 1.** Suppose that $SA_h$ and $ISA_h$ have been computed for some $h > 0$, where $i = SA_h[j]$ is the $j$th suffix in $h$-order and $h$-rank$[i] = ISA_h[i]$. Then, a sort using the integer pairs $(ISA_h[i], ISA_h[i]+h)$ as keys, $i+h \leq n$, computes a $2^h$-order of the suffixes

(Suffixes $i > n-h$ are necessarily already fully ordered.)

The two main prefix-doubling algorithms differ primarily in their application of this observation:

— Algorithm MM does an implicit $2^h$-sort by performing a left-to-right scan of $SA_h$ that induces the $2^h$-rank of $SA_h[j]−h, j = 1, 2, ..., n$;

— Algorithm LS explicitly sorts each $h$-group using the ternary-split quicksort (TSQS) of Bentley and McIlroy [1993].

Algorithm MM employs Observation 1 as follows: If $SA_h$ is scanned left to right (thus, in $h$-order of the suffixes), $j = 1, 2, ..., n$, then the suffixes $i−h = SA_h[j]−h > 0$ are necessarily scanned in $2^h$-order within their respective $h$-groups in $SA_h$. 

source: Puglisi/Smyth/Turpin ACM Computing Surveys ‘07
Induced Sorting

- [Nong/Zhang/Chan DCC’09] **sais**-algorithm:
  - ✓ O(n) in theory
  - ✓ fast in practice
  - ✓ as simple as Kärkkäinen/Sanders DC3
Algorithm sais

- Definition: suffix $T[i,n]$ called
  - **S-type** iff $T[i..n] <_{\text{lex}} T[i+1..n]$ ($T[n,n]=$'\$' always S)
  - **L-type** otherwise

1. Choose sample: leftmost S (call them $S^*$), $|S^*|<1/2n$
2. Sort $S^*$-suffixes by **recursion**
   - on new text formed by sorted $S^*$-substrings
3. Scan $A$ from left to right (say we’re at pos. $i$):
   - if $T[A[i]-1]$ is **L**, write $A[i]-1$ to 1st pos. in bucket
4. like (3), but sorting $S$-suffixes in a right-to-left scan
\[ T = \begin{array}{cccccccccccccc}
0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 & 11 & 12 \\
c & a & b & c & c & b & a & a & a & b & b & a & ~$ \}$ \\
\text{L S*} & \text{S L L L S*} & \text{S S L L L L S*} \\
\end{array} \]
Sorting S*-Substrings

- Same algorithm, but with UNSORTED S*-suffixes

1. Choose sample: leftmost S (call them S*), |S*|<1/2n
2. Put S*-substrings in their buckets (in text order)
3. Scan A from left to right (say we’re at pos. i):
   - if $T[A[i]-1]$ is L, write $A[i]-1$ to 1st pos. in bucket
4. like (3), but sorting S-substrings in a right-to-left scan
Correctness

2 main points:

- S-substrings > L-substrings in same bucket
- order of suffixes in reduced substring \( \preceq \) order in original string

full proof: consult section 3.2 in: