

Trip-Based Public Transit Routing Using Condensed Search Trees

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Agenda



- Introduction
- Trip-Based Public Transit Routing
- Prefix Trees
- Splitting Trees
- Experiments
- Conclusion

Context



- Routing in large, real-world Public Transit Networks
- One-to-one queries
- Pareto-optimal results regarding arrival time and number of transfers
- Earliest arrival and profile/range queries

Terminology



Stops









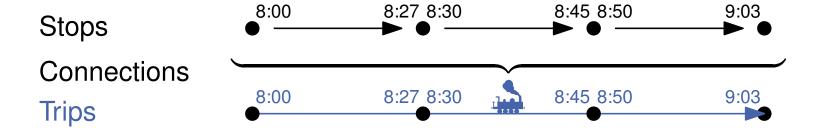
Terminology



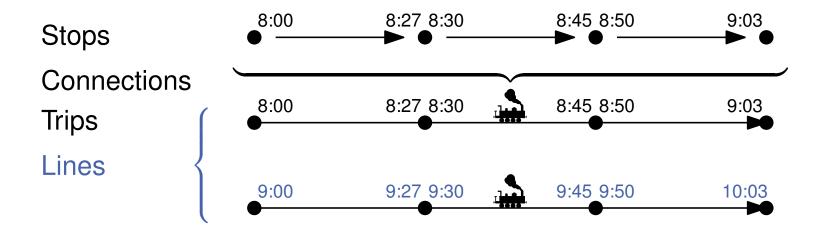


Connections

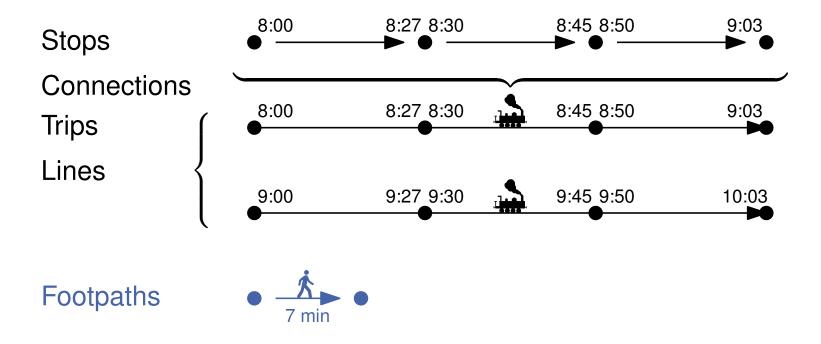




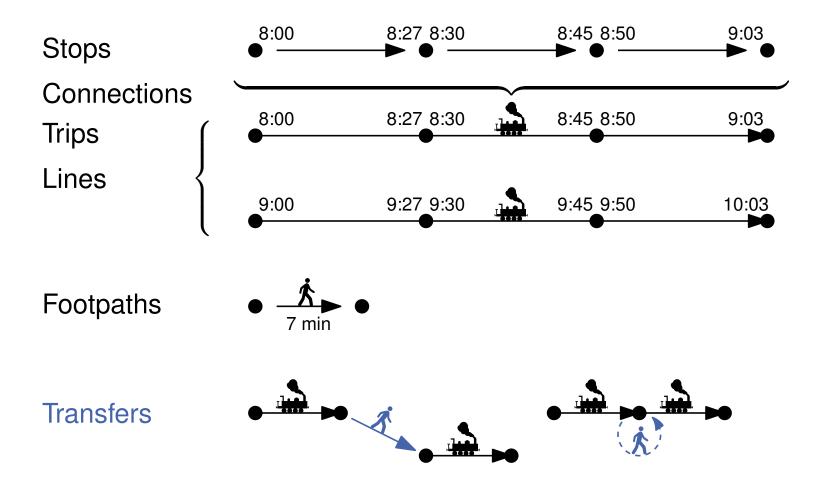




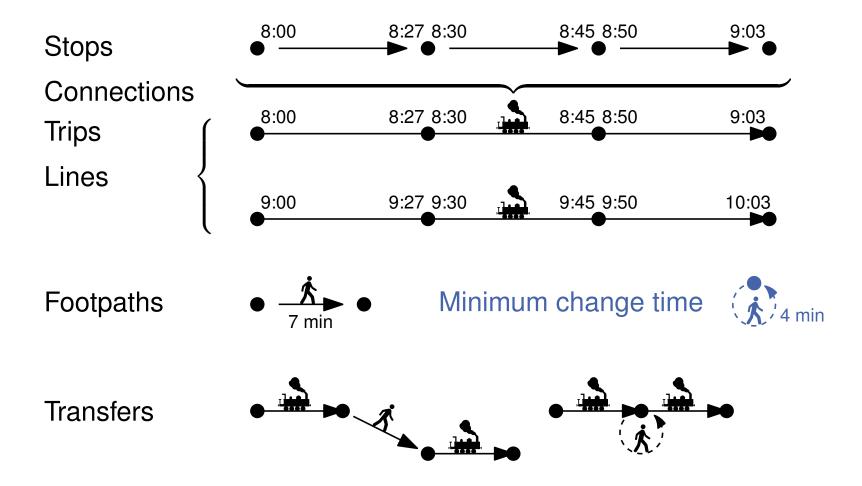




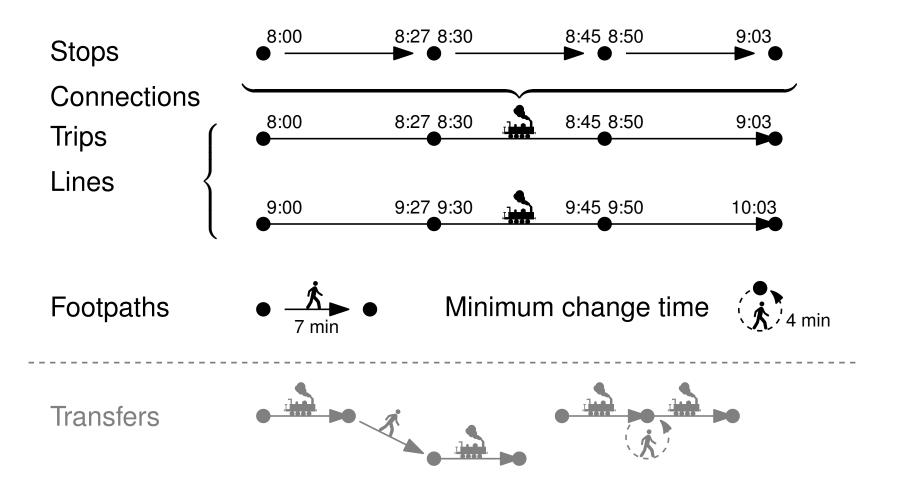








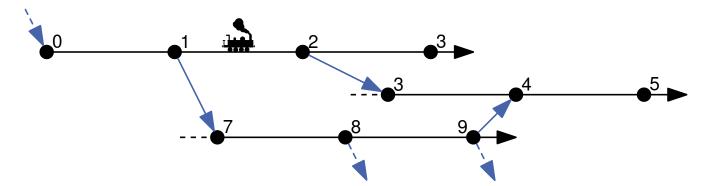




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Intuition

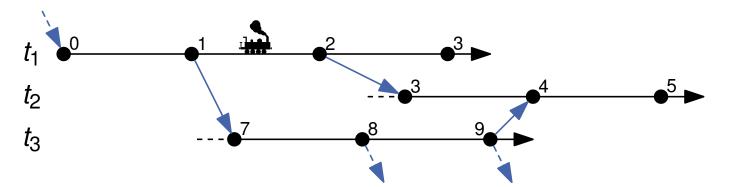
- Move emphasis from stops to trips
- Model transfers between trips explicitly



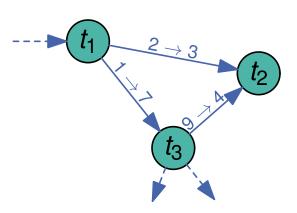
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Intuition

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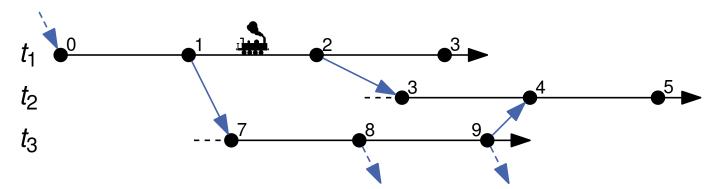
- Query is similar to breadth-first search
- Levels correspond to number of transfers



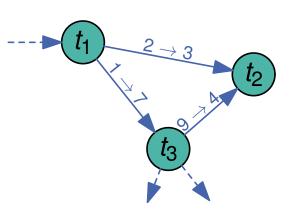
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Intuition

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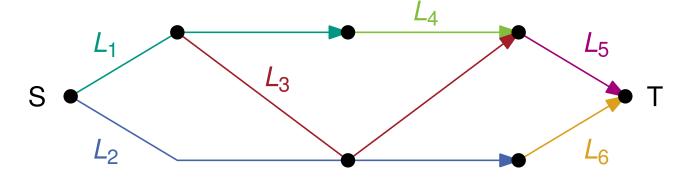


- Query is similar to breadth-first search
- Levels correspond to number of transfers
- Footpaths etc. are handled during preprocessing





Example



<i>L</i> ₁		
8:00	8:30	9:00
:	i	i

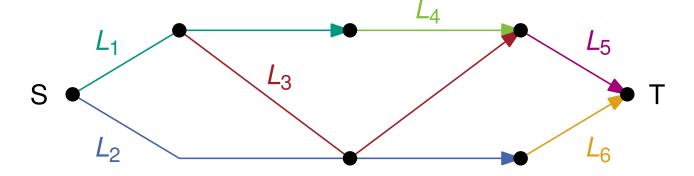
9:30
i

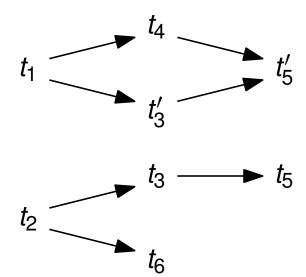
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9:25	9:50
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<i>L</i> ₆	
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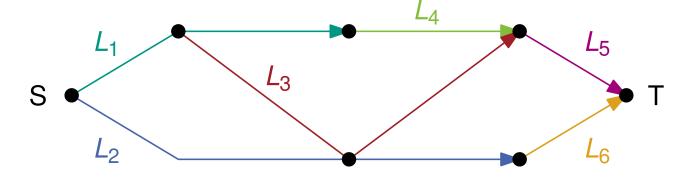
Example

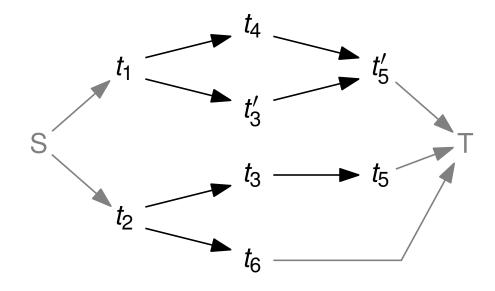




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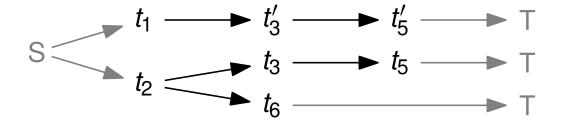
Example





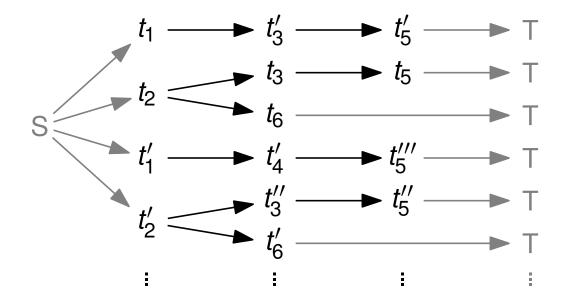
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Search Tree



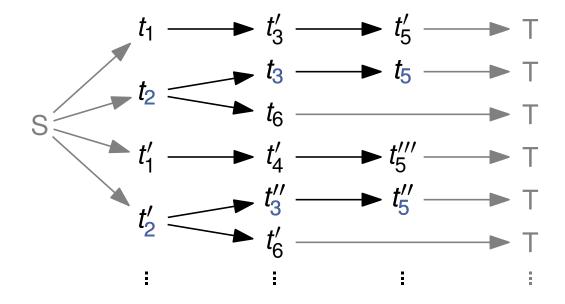
Search Tree





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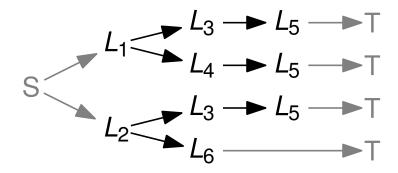
Search Tree



- Repeating patterns
- Can be used for goal-directed search

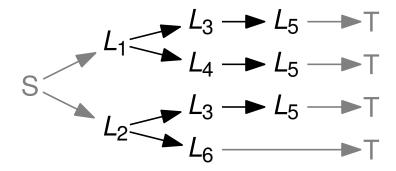
Computation





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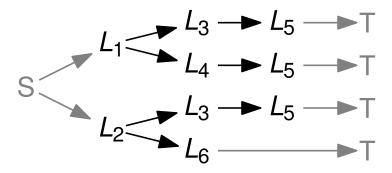
Computation



- Root is source stop, leaves are destination stops
- Each path from S to T forms a sequence of lines that is optimal at some time
- All optimal journeys are covered
- Computed by performing one-to-all profile queries

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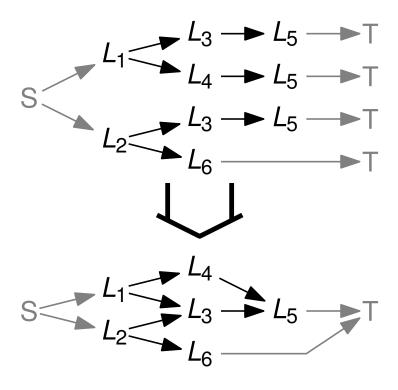
Query



Find all leaves labeled T

Query

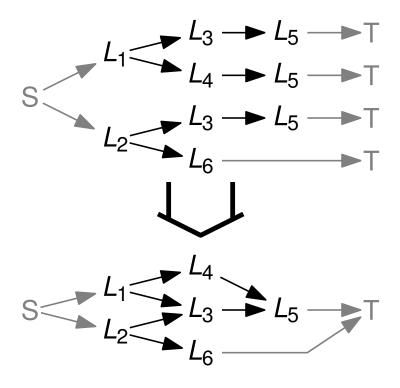




- Find all leaves labeled T
- Collapse paths into graph

Query





- Find all leaves labeled T
- Collapse paths into graph
- Perform query using this graph

Problems

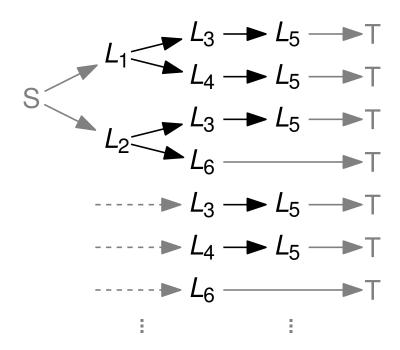


- Large memory consumption
- Each tree covers the whole network

Problems



- Large memory consumption
- Each tree covers the whole network
- A lot of redundancy, especially near the leaves

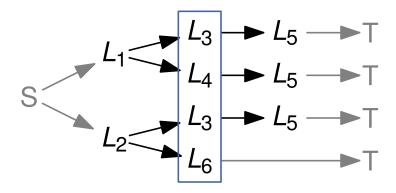




- Solution: Postfix trees
- Can be constructed from prefix trees by cutting off branches

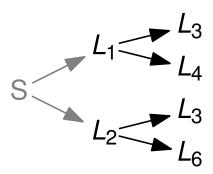
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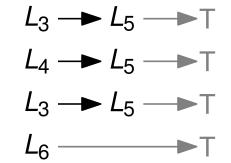
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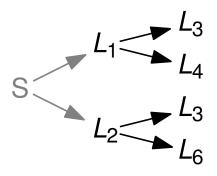
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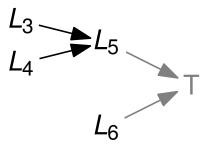




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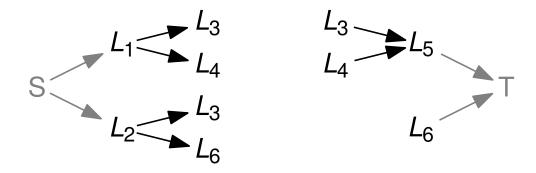
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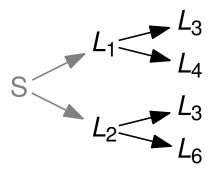


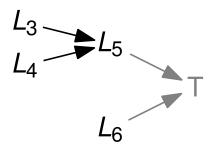
- Need to store two trees for each stop
- But these trees are much smaller

Query



Join trees at matching nodes

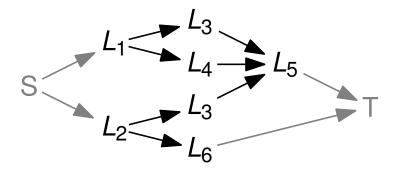




Query



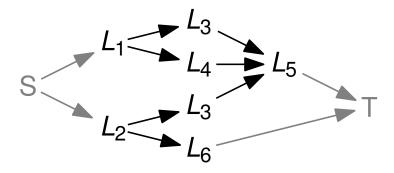
Join trees at matching nodes



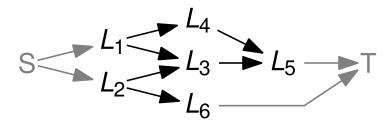
Query



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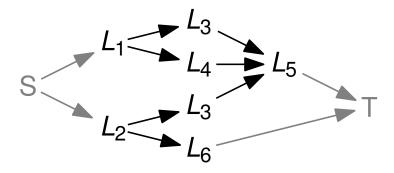
Build query graph as before



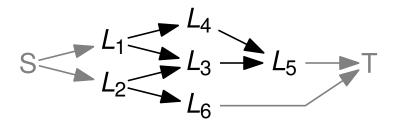
Query



Join trees at matching nodes



Build query graph as before



Run query as before

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Cut Selection

- Cut location is important
- Want to balance prefix and postfix trees, to minimize total size

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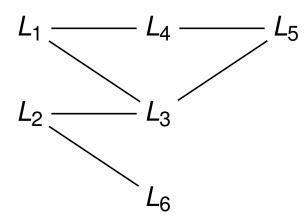
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- Our approach: Select "important" lines
- Intuitively: ICE/TGV > other train > bus

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Cut Selection

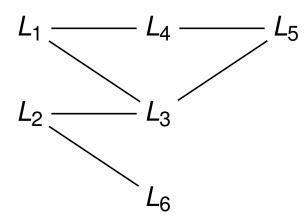
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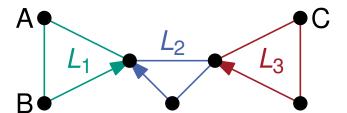


As a result, trees reach out to long-distance lines

There And Back Again



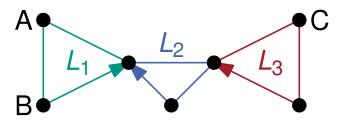
Consider the following graph:



There And Back Again



Consider the following graph:



- L_2 is the most central line
- Prefix and postfix trees for A and B look like this:

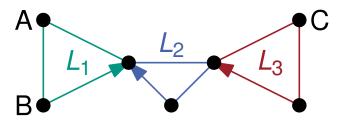
$$A \longrightarrow L_1 \longrightarrow L_2$$

$$L_2 \longrightarrow L_1 \longrightarrow B$$

There And Back Again



Consider the following graph:



- L_2 is the most central line
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Joining them results in this query graph:

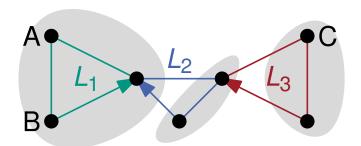
$$A \longrightarrow L_1 \longrightarrow B$$

$$\downarrow L_2$$

There And Back Again



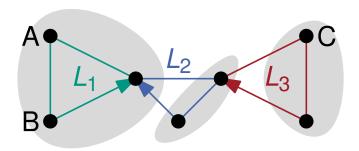
Solution: Partition stops



There And Back Again



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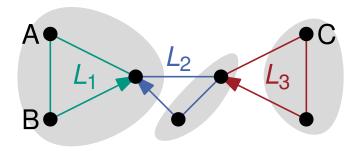
 Use bitmasks to indicate which nodes should be considered for the query graph

$$L_2 \longrightarrow L_1 \longrightarrow B$$

There And Back Again



Solution: Partition stops



 Use bitmasks to indicate which nodes should be considered for the query graph

$$A \longrightarrow L_1 \longrightarrow L_2$$
100 011

$$L_2 \longrightarrow L_1 \longrightarrow B$$

Result: Fewer false positives

$$A \longrightarrow L_1 \longrightarrow B$$





Instance	Stops	Conn.	Trips	Lines	Footp.	Transfers
Germany	296.6 k	27,062 k	1,432 k	192.9 k	102.8 k	84,953 k
Sweden	50.7 k	6,054 k	261 k	17.6 k	$0.8\mathrm{k}$	16,455 k
Switzerland	27.8 k	4,650 k	611 k	14.4 k	34.3 k	12,626 k
London	20.8 k	4,991 k	129 k	2.2 k	27.6 k	15,883 k
Madrid	4.6 k	5,280 k	190 k	1.4 k	1.4 k	9,256 k





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	p. prefix	seq.	par.	speed	avg.#	mem.
Instance	tree [ms]	[h:m]	[h:m]	up	of nodes	[GB]
Germany	2,143.6	(231:16)	13:48	(16.8 x)	6,131	23.2
Sweden	166.7	4:33	0:18	15.2 x	2,433	1.6
Switzerland	209.3	3:18	0:12	16.5 x	4,315	1.6
London	1,368.1	15:19	0:42	21.9 x	20,390	6.0
Madrid	497.3	1:22	0:04	17.0 x	32,293	2.0

^{*} Quad 8-core Intel Xeon E5-4640, 2.4 GHz, 512 GB DDR3-1600, 64 threads





		Query graph	Query graph	EA	profile
Instance	Var.	size [N+E]	time [µs]	[µs]	[μs]
Germany	TB			30,856	192,952
Sweden	TB			2,760	16,532
Switzerland	TB			1,780	18,104
London	TB			1,374	96,114
Madrid	TB	_	_	711	54,118
Germany	PT	41 + 58	994.4	63.3	155.0
Sweden	PT	23 + 32	24.6	40.4	88.6
Switzerland	PT	38 + 59	34.0	45.8	155.9
London	PT	91 + 196	138.2	101.1	2,786.6
Madrid	PT	150 + 407	306.9	81.7	6,913.8
Germany	ST	124 + 232	81.1	75.0	430.5
Sweden	ST	66 + 122	32.5	27.2	207.1
Switzerland	ST	118 + 233	76.1	32.7	327.6
London	ST	331 + 1,242	1,583.3	141.4	14,545.4
Madrid	ST	456 + 2,073	11,822.9	165.8	28,919.0





algorithm	instance	stops [10 ³]	conn. [10 ⁶]	AS A	O'O'III	mem. [GB]	pre. [h]	query [μs]
CSA	Germany	252.4	46.2	0	0			298.6 k
ACSA	Germany	252.4	46.2	0	0	n/a	0.2	8.7 k
TP	Germany	248.4	13.9	•	0	140.0	372.0	300.0
Sc-TP	Germany	250.0	15.0	•	0	1.2	16.5	32.0 k
TB	Germany	296.6	27.1	•	0	23.2	231.3	156.1
TTL	Sweden	51.4	n/a	0	0	\approx 0.5	0.2	\approx 10.0
PTL	Sweden	51.1	12.7	•	0	12.3	36.2	27.6
TB	Sweden	50.7	6.1	•	0	1.6	3.8	59.7
PTL	Switzerland	27.1	23.7	•	0	12.7	61.6	21.7
TB	Switzerland	27.8	4.7	•	0	1.6	2.7	108.8
CSA	London	20.8	4.9	0	0			1.8 k
PTL	London	20.8	5.1	•	0	26.2	49.3	30.0
TB	London	20.8	5.0	•	0	6.0	11.6	1.7 k
TTL	Madrid	4.6	n/a	0	0	≈ 0.4	0.1	\approx 30.0
PTL	Madrid	4.7	4.5	•	0	9.9	10.9	64.3
TP	Madrid	4.6	4.8	•	0	n/a	185.0	3.1 k
ТВ	Madrid	4.6	5.3	•	0	2.0	1.1	12.0 k





algorithm	instance	stops [10 ³]	conn. [10 ⁶]	ASI OF	O O ST.	mem. [GB]	pre. [h]	query [μs]
ACSA TP TB	Germany Germany Germany	252.4 248.4 296.6	46.2 13.9 27.1	•	•	n/a 140.0 23.2	0.2 372.0 231.3	171.0 k 5.0 k 511.6
PTL TB	Sweden Sweden	51.1 50.7	12.7 6.1	•	•	0.7 1.6	0.5 3.8	12.1 239.6
PTL TB	Switzerland Switzerland	27.1 27.8	23.7 4.7	•	•	0.7 1.6	0.7 2.7	24.5 403.7
PTL CSA TB	London London London	20.8 20.8 20.8	5.1 4.9 5.0	•	•	1.3	0.9	74.3 466.0 k 16.1 k
PTL TB	Madrid Madrid	4.7 4.6	4.5 5.3	•	•	0.4 2.0	0.4 1.1	111.9 40.7 k

Conclusion



- Speed-up technique for trip-based public transit routing
- Exploits regularities in timetables
- Fast preprocessing of country-sized networks
- Pareto-optimal bi-criteria 24 h profile queries on microsecond scale

Conclusion



- Speed-up technique for trip-based public transit routing
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Future Work

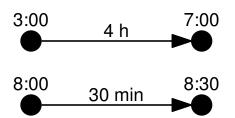
- Improve preprocessing time (and space) for very large networks
- Improve performance on metropolitan networks
- Better scaling to larger networks and longer timeframes

Introduction

Public Transit Routing



Inherently time-dependent:
 Travel times depend on departure time

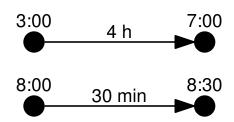


Introduction

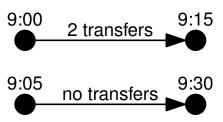
Public Transit Routing



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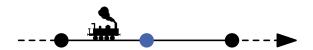


- Multiple natural problem variants
 - Earliest arrival queries
 - Profile (range) queries
 - Multi-criteria queries (e.g., number of transfers taken)



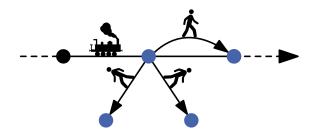
Preprocessing





Preprocessing

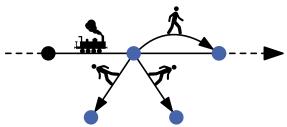
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Preprocessing



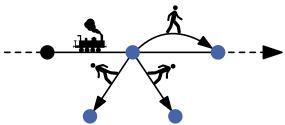
Compute transfers between trips



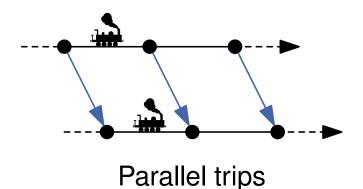
For each line, find the first reachable trip $(arrival time + footpath length \le departure time)$

Preprocessing



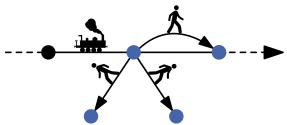


- For each line, find the first reachable trip (arrival time + footpath length \leq departure time)
- Huge number of transfers, not all of which are useful

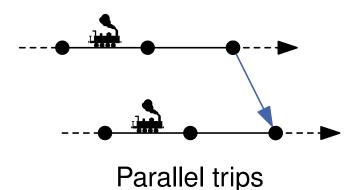


Preprocessing



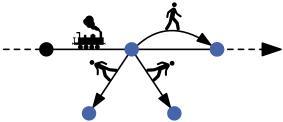


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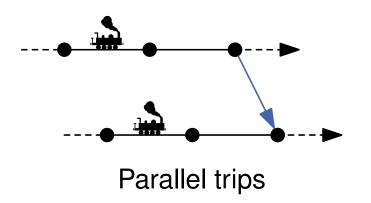


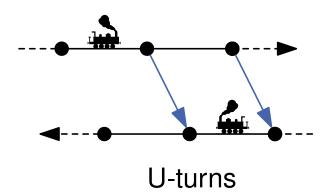
Preprocessing





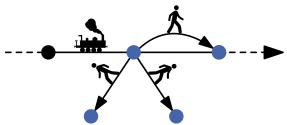
- For each line, find the first reachable trip (arrival time + footpath length ≤ departure time)
- Huge number of transfers, not all of which are useful



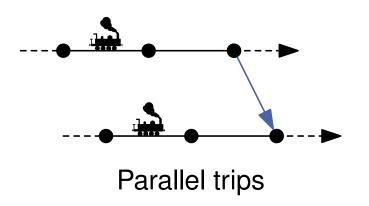


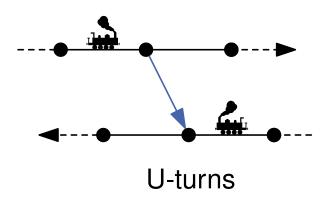
Preprocessing





- For each line, find the first reachable trip (arrival time + footpath length ≤ departure time)
- Huge number of transfers, not all of which are useful





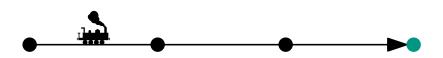


Preprocessing

Reduce number of transfers by eliminating redundant ones

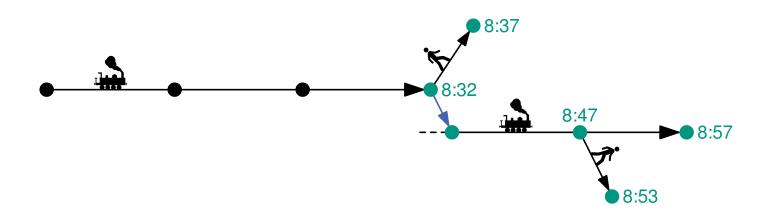
Karlsruhe Institute of Technology

- Reduce number of transfers by eliminating redundant ones
- Process trips backwards



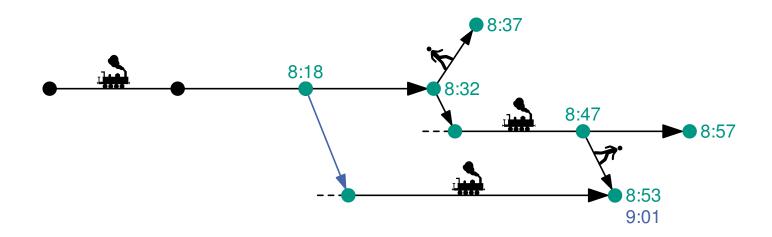


- Reduce number of transfers by eliminating redundant ones
- Process trips backwards
- Keep track of which stops can be reached at what time



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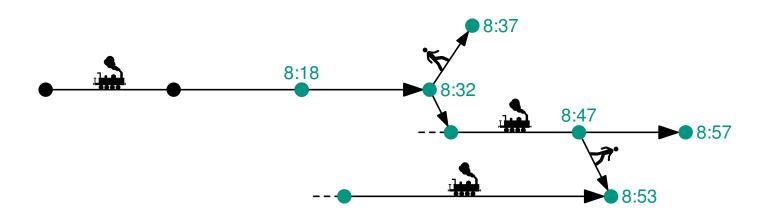
- Reduce number of transfers by eliminating redundant ones
- Process trips backwards
- Keep track of which stops can be reached at what time
- Evaluate transfers by checking if they improve arrival times





- Reduce number of transfers by eliminating redundant ones
- Process trips backwards

- Keep track of which stops can be reached at what time
- Evaluate transfers by checking if they improve arrival times
- Removes up to 90% of original transfers



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Query

Input: Source stop, target stop, departure time

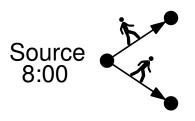
Source 8:00

Target

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Query

- Input: Source stop, target stop, departure time
- Identify trips reachable from the source



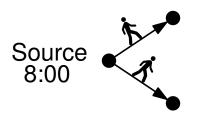
Target

	don	lina	trin	indov
	dep.	ime	trip	index
	8:00	2	15	8
	8:03	4	56	0
	8:07	11	456	31
	9:00	110	3256	6
•				

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Query

- Input: Source stop, target stop, departure time
- Identify trips reachable from the source
- Identify lines reaching the target





dep.	line	trip	index
8:00	2	15	8
8:03	4	56	0
8:07	11	456	31
9:00	110	3256	6

line	index	footpath
3	8	
8	17	4 min
27	3	4 min

Query

Queue trips and mark as reached



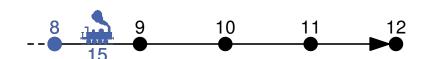
	Queu	е	Reach	ned
# tr.	trip	range		_
0	15	8–12	15 16	8
0	56	0-14	17	8
0	456	31–78		•
0	3256	6–45	56 57	0
			450	~ 4
			456 457	31 31

dep.	line	trip	index
8:00	2	15	8
8:03	4	56	0
8:07	11	456	31
9:00	110	3256	6

•	line	index	footpath
•	3	8	
	8	17	4 min
	27	3	4 min

Query

- Queue trips and mark as reached
- Process queue



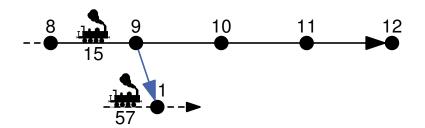


	Queu	е	Reac	hed
# tr.	trip	range	15	. 8
0	15	8–12	16	
0	56	0–14	17	8 8
0	456	31–78	F.C.	
0	3256	6–45	56 57	0
			456 457	31 31
			3256	6

line	index	footpath
3	8	
8	17	4 min
27	3	4 min

Query

- Queue trips and mark as reached
- Process queue
- Examine transfers



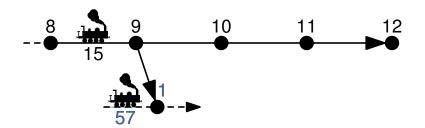


Reache	е	Queu	
1	range	trip	# tr.
15 16	8–12	15	0
17	0–14	56	0
	31–78	456	0
56 57	6–45	3256	0
456 3 457 3			
3256			

•	line	index	footpath
•	3	8	
	8	17	4 min
	27	3	4 min

Query

- Queue trips and mark as reached
- Process queue
- Examine transfers
 - Compare against label





Panchad

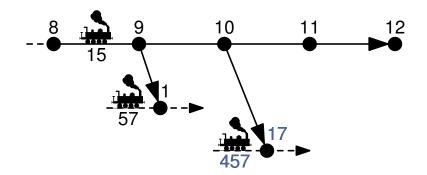
	Queu	е		neaci	ieu
# tr.	trip	range	_	15	 8
0	15	8–12		16	8
0	56	0–14		17	8
0	456	31–78			
0	3256	6–45		56	0
U	0230	0 73		57	0
				456	31
				457	31

 Ω

•	line	index	footpath
•	3	8	
	8	17	4 min
	27	3	4 min

Query

- Queue trips and mark as reached
- Process queue
- Examine transfers
 - Compare against label



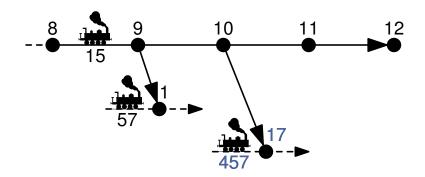


	Queu	е	Reach	ned
# tr.	trip	range	 15	8
0	15	8–12	16	8
0	56	0–14	17	8
0	456	31–78	 F0	•
0	3256	6–45	56 57	0
			456	31
			457	31

line	index	footpath
3	8	
8	17	4 min
27	3	4 min

Query

- Queue trips and mark as reached
- Process queue
- Examine transfers
 - Compare against label
 - Queue newly reached trips





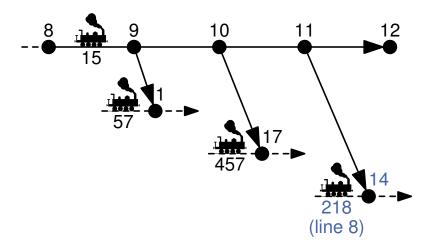
	Queu	е		Reacl	ned
# tr.	trip	range	•	15	. 0
0	15	8–12		16	8 8
0	56	0–14		17	8 8
0	456	31–78		· · ·	•
0	3256	6–45		56 57	0
1	457	17–31			U
				456	31
				457	17
				458	17

line	index	footpath
3	8	
8	17	4 min
27	3	4 min





- Queue trips and mark as reached
- Process queue
- Examine transfers
 - Compare against label
 - Queue newly reached trips
 - Output a journey if target is reached



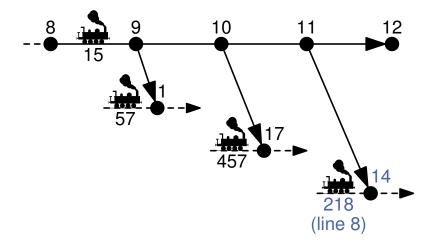
	Queu	е		Reach	ned
# tr.	trip	range	-	1.5	
0	15	8–12		15 16	8 8
0	56	0–14		17	8
0	456	31–78			•
0	3256	6–45		56 57	0
1	457	17–31			U
1	218	14–23		218	14
				456 457 458	31 17 17

line	index	footpath
3	8	
8	17	4 min
27	3	4 min

Query



- Queue trips and mark as reached
- Process queue
- Examine transfers
 - Compare against label
 - Queue newly reached trips
 - Output a journey if target is reached



ached	Reac	Queue		
 E O	15	range	trip	# tr.
	15 16	8–12	15	0
6 8 7 8	17	0–14	56	0
	 	31–78	456	0
	56 57	6–45	3256	0
		17–31	457	1
8 14	218	14–23	218	1
 'C 01	156			

line	index	footpath
3	8	
8	17	4 min
27	3	4 min

 $arrival_time(218, 17) = 9:24 \Longrightarrow Arrival at 9:28 after 1 transfer$





Danchad

- Queue trips and mark as reached
- Process queue
- Examine transfers
 - Compare against label
 - Queue newly reached trips
 - Output a journey if target is reached
- Continue until queue is empty

ieu	neaci	e	Queue		
8	15	range	trip	# tr.	
8	16	8–12	15	0	
8	17	0–14	56	0	
^		31–78	456	0	
0	56 57	6–45	3256	0	
·	• • •	17–31	457	1	
14	218	14–23	218	1	
31 17 17	456 457 458				

 Ω

line	index	footpath
3	8	
8	17	4 min
27	3	4 min

Arrival at 9:28 after 1 transfer





- Queue trips and mark as reached
- Process queue
- Examine transfers
 - Compare against label
 - Queue newly reached trips
 - Output a journey if target is reached
- Continue until queue is empty
- Skip trips that cannot improve the currently best arrival time

departure_time(1302, 8) = 9:32 > 9:28

Queue			Reacl	ned
# tr.	trip	range	15	. 0
1	1302	8–45	16	8 8
1	2871	3–11	17	8
2	512	0–19	· · ·	•
2	1523	19–88	56 57	0
2	43	13–15		·
2	44	4–53	218	14
			456 457 458	31 17 17
			3256	6

line	index	footpath
	IIIGOX	Ισστρατίτ
3	8	
8	17	4 min
27	3	4 min

Arrival at 9:28 after 1 transfer

Query



- Process queue
- Examine transfers
 - Compare against label
 - Queue newly reached trips
 - Output a journey if target is reached
- Continue until queue is empty
- Skip trips that cannot improve the currently best arrival time



Dagahad

nea	Reaci	Queue		
. 0	15	range	trip	# tr.
8 8 8	16	3–11	2871	1
8	17	0–19	512	2
	 50	19–88	1523	2
0	56 57	13–15	43	2
Ū		4-53	44	2
14	218			
31 17 17	456 457 458			

line	index	footpath
3	8	
8	17	4 min
27	3	4 min

Arrival at 9:28 after 1 transfer